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Solutions Manual to Walter  
Rudin's *Principles of  
Mathematical Analysis*

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## Chapter 1

# The Real and Complex Number Systems

**Exercise 1.1** If  $r$  is rational ( $r \neq 0$ ) and  $x$  is irrational, prove that  $r + x$  and  $rx$  are irrational.

*Solution.* If  $r$  and  $r + x$  were both rational, then  $x = r + x - r$  would also be rational. Similarly if  $rx$  were rational, then  $x = \frac{rx}{r}$  would also be rational.

**Exercise 1.2** Prove that there is no rational number whose square is 12.

*First Solution.* Since  $\sqrt{12} = 2\sqrt{3}$ , we can invoke the previous problem and prove that  $\sqrt{3}$  is irrational. If  $m$  and  $n$  are integers having no common factor and such that  $m^2 = 3n^2$ , then  $m$  is divisible by 3 (since if  $m^2$  is divisible by 3, so is  $m$ ). Let  $m = 3k$ . Then  $m^2 = 9k^2$ , and we have  $3k^2 = n^2$ . It then follows that  $n$  is also divisible by 3 contradicting the assumption that  $m$  and  $n$  have no common factor.

*Second Solution.* Suppose  $m^2 = 12n^2$ , where  $m$  and  $n$  have no common factor. It follows that  $m$  must be even, and therefore  $n$  must be odd. Let  $m = 2r$ . Then we have  $r^2 = 3n^2$ , so that  $r$  is also odd. Let  $r = 2s + 1$  and  $n = 2t + 1$ . Then

$$4s^2 + 4s + 1 = 3(4t^2 + 4t + 1) = 12t^2 + 12t + 3,$$

so that

$$4(s^2 + s - 3t^2 - 3t) = 2.$$

But this is absurd, since 2 cannot be a multiple of 4.

**Exercise 1.3** Prove Proposition 1.15, i.e., prove the following statements:

- (a) If  $x \neq 0$  and  $xy = xz$ , then  $y = z$ .
- (b) If  $x \neq 0$  and  $xy = x$ , then  $y = 1$ .
- (c) If  $x \neq 0$  and  $xy = 1$ , then  $y = 1/x$ .
- (d) If  $x \neq 0$ , then  $1/(1/x) = x$ .

*Solution.* (a) Suppose  $x \neq 0$  and  $xy = xz$ . By Axiom (M5) there exists an element  $1/x$  such that  $(1/x)x = 1$ . By (M3) and (M4) we have  $(1/x)(xy) = ((1/x)x)y = 1y = y$ , and similarly  $(1/x)(xz) = z$ . Hence  $y = z$ .

- (b) Apply (a) with  $z = 1$ .
- (c) Apply (a) with  $z = 1/x$ .
- (d) Apply (a) with  $x$  replaced by  $1/x$ ,  $y = 1/(1/x)$ , and  $z = x$ .

**Exercise 1.4** Let  $E$  be a nonempty subset of an ordered set; suppose  $\alpha$  is a lower bound of  $E$ , and  $\beta$  is an upper bound of  $E$ . Prove that  $\alpha \leq \beta$ .

*Solution.* Since  $E$  is nonempty, there exists  $x \in E$ . Then by definition of lower and upper bounds we have  $\alpha \leq x \leq \beta$ , and hence by property *ii* in the definition of an ordering, we have  $\alpha < \beta$  unless  $\alpha = x = \beta$ .

**Exercise 1.5** Let  $A$  be a nonempty set of real numbers which is bounded below. Let  $-A$  be the set of all numbers  $-x$ , where  $x \in A$ . Prove that

$$\inf A = -\sup(-A).$$

*Solution:* We need to prove that  $-\sup(-A)$  is the greatest lower bound of  $A$ . For brevity, let  $\alpha = -\sup(-A)$ . We need to show that  $\alpha \leq x$  for all  $x \in A$  and  $\alpha \geq \beta$  if  $\beta$  is any lower bound of  $A$ .

Suppose  $x \in A$ . Then,  $-x \in -A$ , and, hence  $-x \leq \sup(-A)$ . It follows that  $x \geq -\sup(-A)$ , i.e.,  $\alpha \leq x$ . Thus  $\alpha$  is a lower bound of  $A$ .

Now let  $\beta$  be any lower bound of  $A$ . This means  $\beta \leq x$  for all  $x$  in  $A$ . Hence  $-\beta \leq -x$  for all  $x \in A$ , which says  $-\beta \leq y$  for all  $y \in -A$ . This means  $-\beta$  is an upper bound of  $-A$ . Hence  $-\beta \geq \sup(-A)$  by definition of  $\sup$ , i.e.,  $\beta \leq -\sup(-A)$ , and so  $-\sup(-A)$  is the greatest lower bound of  $A$ .

**Exercise 1.6** Fix  $b > 1$ .

- (a) If  $m, n, p, q$  are integers,  $n > 0, q > 0$ , and  $r = m/n = p/q$ , prove that

$$(b^m)^{1/n} = (b^p)^{1/q}.$$

Hence it makes sense to define  $b^r = (b^m)^{1/n}$ .

- (b) Prove that  $b^{r+s} = b^r b^s$  if  $r$  and  $s$  are rational.

(c) If  $x$  is real, define  $B(x)$  to be the set of all numbers  $b^t$ , where  $t$  is rational and  $t \leq x$ . Prove that

$$b^r = \sup B(r)$$

when  $r$  is rational. Hence it makes sense to define

$$b^x = \sup B(x)$$

for every real  $x$ .

(d) Prove that  $b^{x+y} = b^x b^y$  for all real  $x$  and  $y$ .

*Solution.* (a) Let  $k = mq = np$ . Since there is only one positive real number  $c$  such that  $c^{nq} = b^k$  (Theorem 1.21), if we prove that both  $(b^m)^{1/n}$  and  $(b^p)^{1/q}$  have this property, it will follow that they are equal. The proof is then a routine computation:  $((b^m)^{1/n})^{nq} = (b^m)^q = b^{mq} = b^k$ , and similarly for  $(b^p)^{1/q}$ .

(b) Let  $r = \frac{m}{n}$  and  $s = \frac{v}{w}$ . Then  $r + s = \frac{mw+vn}{nw}$ , and

$$b^{r+s} = (b^{mw+vn})^{1/nw} = ((b^{mw}b^{vn}))^{1/nw},$$

by the laws of exponents for integer exponents. By the corollary to Theorem 1.21 we then have

$$b^{r+s} = (b^{mw})^{1/nw} (b^{vn})^{1/nw} = b^r b^s,$$

where the last equality follows from part (a).

(c) It will simplify things later on if we amend the definition of  $B(x)$  slightly, by defining it as  $\{b^t : t \text{ rational}, t < x\}$ . It is then slightly more difficult to prove that  $b^r = \sup B(r)$  if  $r$  is rational, but the technique of Problem 7 comes to our rescue. Here is how: It is obvious that  $b^r$  is an upper bound of  $B(r)$ . We need to show that it is the least upper bound. The inequality  $b^{1/n} < t$  if  $n > (b-1)/(t-1)$  is proved just as in Problem 7 below. It follows that if  $0 < x < b^r$ , there exists an integer  $n$  with  $b^{1/n} < b^r/x$ , i.e.,  $x < b^{r-1/n} \in B(r)$ . Hence  $x$  is not an upper bound of  $B(r)$ , and so  $b^r$  is the least upper bound.

(d) By definition  $b^{x+y} = \sup B(x+y)$ , where  $B(x+y)$  is the set of all numbers  $b^t$  with  $t$  rational and  $t < x+y$ . Now any rational number  $t$  that is less than  $x+y$  can be written as  $r+s$ , where  $r$  and  $s$  are rational,  $r < x$ , and  $s < y$ . To do this, let  $r$  be any rational number satisfying  $t-y < r < x$ , and let  $s = t-r$ . Conversely any pair of rational numbers  $r, s$  with  $r < x$ ,  $s < y$  gives a rational sum  $t = r+s < x+y$ . Hence  $B(x+y)$  can be described as the set of all numbers  $b^r b^s$  with  $r < x$ ,  $s < y$ , and  $r$  and  $s$  rational, i.e.,  $B(x+y)$  is the set of all products  $uv$ , where  $u \in B(x)$  and  $v \in B(y)$ .

Since any such product is less than  $\sup B(x) \sup B(y)$ , we see that the number  $M = \sup B(x) \sup B(y)$  is an upper bound for  $B(x+y)$ . On the other hand, suppose  $0 < c < \sup B(x) \sup B(y)$ . Then  $c/(\sup B(x)) < \sup B(y)$ . Let  $m = (1/2)(c/(\sup B(x)) + \sup B(y))$ . Then  $c/\sup B(x) < m < \sup B(y)$ , and there exist  $u \in B(x)$ ,  $v \in B(y)$  such that  $c/m < u$  and  $m < v$ . Hence we have

$c = (c/m)m < uv \in B(x + y)$ , and so  $c$  is not an upper bound for  $B(x + y)$ . It follows that  $\sup B(x) \sup B(y)$  is the least upper bound of  $B(x + y)$ , i.e.,

$$b^{x+y} = b^x b^y,$$

as required.

**Exercise 1.7** Fix  $b > 1$ ,  $y > 0$ , and prove that there is a unique real  $x$  such that  $b^x = y$ , by completing the following outline. (This  $x$  is called the *logarithm of  $y$  to the base  $b$* .)

(a) For any positive integer  $n$ ,  $b^n - 1 \geq n(b - 1)$ .

(b) Hence  $b - 1 \geq n(b^{1/n} - 1)$ .

(c) If  $t > 1$  and  $n > (b - 1)/(t - 1)$ , then  $b^{1/n} < t$ .

(d) If  $w$  is such that  $b^w < y$ , then  $b^{w+(1/n)} < y$  for sufficiently large  $n$ ; to see this apply part (c) with  $t = y \cdot b^{-w}$ .

(e) If  $b^w > y$ , then  $b^{w-(1/n)} > y$  for sufficiently large  $n$ .

(f) Let  $A$  be the set of all  $w$  such that  $b^w < y$ , and show that  $x = \sup A$  satisfies  $b^x = y$ .

(g) Prove that this  $x$  is unique.

*Solution.* (a) The inequality  $b^n - 1 \geq n(b - 1)$  is equality if  $n = 1$ . Then, by induction  $b^{n+1} - 1 = b^{n+1} - b + (b - 1) = b(b^n - 1) + (b - 1) \geq bn(b - 1) + (b - 1) = (bn + 1)(b - 1) \geq (n + 1)(b - 1)$ .

(b) Replace  $b$  by  $b^{1/n}$  in part (a).

(c) The inequality  $n > (b - 1)/(t - 1)$  can be rewritten as  $n(t - 1) > (b - 1)$ , and since  $b - 1 \geq n(b^{1/n} - 1)$ , we have  $n(t - 1) > n(b^{1/n} - 1)$ , which implies  $t > b^{1/n}$ .

(d) The application of part (c) with  $t = y \cdot b^{-w} > 1$  is immediate.

(e) The application of part (c) with  $t = b^w \cdot (1/y)$  yields the result, as in part (d) above.

(f) There are only three possibilities for the number  $x = \sup A$ : 1)  $b^x < y$ ; 2)  $b^x > y$ ; 3)  $b^x = y$ . The first assumption, by part (d), implies that  $x + (1/n) \in A$  for large  $n$ , contradicting the assumption that  $x$  is an upper bound for  $A$ . The second, by part (e), implies that  $x - (1/n)$  is an upper bound for  $A$  if  $n$  is large, contradicting the assumption that  $x$  is the smallest upper bound. Hence the only remaining possibility is that  $b^x = y$ .

(g) Suppose  $z \neq x$ , say  $z > x$ . Then  $b^z = b^{x+(z-x)} = b^x b^{z-x} > b^x = y$ . Hence  $x$  is unique. (It is easy to see that  $b^w > 1$  if  $w > 0$ , since there is a positive rational number  $r = \frac{m}{n}$  with  $0 < r < w$ , and  $b^r = (b^m)^{1/n}$ . Then  $b^m > 1$  since  $b > 1$ , and  $(b^m)^{1/n} > 1$  since  $1^n = 1 < b^m$ .)

**Exercise 1.8** Prove that no order can be defined in the complex field that turns it into an ordered field. *Hint:*  $-1$  is a square.

*Solution.* By Part (a) of Proposition 1.18, either  $i$  or  $-i$  must be positive. Hence  $-1 = i^2 = (-i)^2$  must be positive. But then  $1 = (-1)^2$ , must also be positive, and this contradicts Part (a) of Proposition 1.18, since  $1$  and  $-1$  cannot both be positive.

**Exercise 1.9** Suppose  $z = a + bi$ ,  $w = c + di$ . Define  $z < w$  if  $a < c$ , and also if  $a = c$  but  $b < d$ . Prove that this turns the set of all complex numbers into an ordered set. (This type of order relation is called a *dictionary order*, or *lexicographic order*, for obvious reasons.) Does this ordered set have the least upper bound property?

*Solution.* We need to show that either  $z < w$  or  $z = w$ , or  $w < z$ . Now since the *real* numbers are ordered, we have  $a < c$  or  $a = c$ , or  $c < a$ . In the first case  $z < w$ ; in the third case  $w < z$ . Now consider the second case. We must have  $b < d$  or  $b = d$  or  $d < b$ . In the first of these cases  $z < w$ , in the third case  $w < z$ , and in the second case  $z = w$ .

We also need to show that if  $z < w$  and  $w < u$ , then  $z < u$ . Let  $u = e + fi$ . Since  $z < w$ , we have either  $a < c$  or  $a = c$  and  $b < d$ . Since  $w < u$  we have either  $c < e$  or  $c = e$  and  $d < f$ . Hence there are four possible cases:

Case 1:  $a < c$  and  $c < e$ . Then  $a < e$  and so  $z < u$ , as required.

Case 2:  $a < c$  and  $c = e$  and  $d < f$ . Again  $a < e$ , and  $z < u$ .

Case 3:  $a = c$  and  $b < d$  and  $c < e$ . Once again  $a < e$  and so  $z < u$ .

Case 4:  $a = c$  and  $b < d$  and  $c = e$ , and  $d < f$ . Then  $a = e$  and  $b < f$ , and so  $z < u$ .

**Exercise 1.10** Suppose  $z = a + bi$ ,  $w = u + iv$ , and

$$a = \left( \frac{|w| + u}{2} \right)^{1/2}, \quad b = \left( \frac{|w| - u}{2} \right)^{1/2}.$$

Prove that  $z^2 = w$  if  $v \geq 0$  and that  $(\bar{z})^2 = w$  if  $v \leq 0$ . Conclude that every complex number (with one exception) has two complex square roots.

*Solution.*

$$z^2 = (a + bi)^2 = (a^2 - b^2) + 2abi.$$

Now

$$a^2 - b^2 = \frac{|w| + u}{2} - \frac{|w| - u}{2} = u,$$

and, since  $(xy)^{1/2} = x^{1/2}y^{1/2}$ ,

$$2ab = 2 \left( \frac{|w| + u}{2} \frac{|w| - u}{2} \right)^{1/2} = 2 \left( \frac{|w|^2 - u^2}{4} \right)^{1/2}.$$

Hence

$$2ab = 2\left(\left(\frac{v}{2}\right)^2\right)^{1/2}$$

Now  $(x^2)^{1/2} = x$  if  $x \geq 0$  and  $(x^2)^{1/2} = -x$  if  $x \leq 0$ . We conclude that  $2ab = v$  if  $v \geq 0$  and  $2ab = -v$  if  $v \leq 0$ . Hence  $z^2 = w$  if  $v \geq 0$ . Replacing  $b$  by  $-b$ , we find that  $(\bar{z})^2 = w$  if  $v \leq 0$ .

Hence every non-zero complex number has (at least) two complex square roots.

**Exercise 1.11** If  $z$  is a complex number, prove that there exists an  $r \geq 0$  and a complex number  $w$  with  $|w| = 1$  such that  $z = rw$ . Are  $w$  and  $r$  always uniquely determined by  $z$ ?

*Solution.* If  $z = 0$ , we take  $r = 0$ ,  $w = 1$ . (In this case  $w$  is not unique.) Otherwise we take  $r = |z|$  and  $w = z/|z|$ , and these choices are unique, since if  $z = rw$ , we must have  $r = r|w| = |rw| = |z|$ ,  $z/r$ .

**Exercise 1.12** If  $z_1, \dots, z_n$  are complex, prove that

$$|z_1 + z_2 + \cdots + z_n| \leq |z_1| + |z_2| + \cdots + |z_n|.$$

*Solution.* The case  $n = 2$  is Part (e) of Theorem 1.33. We can then apply this result and induction on  $n$  to get

$$\begin{aligned} |z_1 + z_2 + \cdots + z_n| &= |(z_1 + z_2 + \cdots + z_{n-1}) + z_n| \\ &\leq |z_1 + z_2 + \cdots + z_{n-1}| + |z_n| \\ &\leq |z_1| + |z_2| + \cdots + |z_{n-1}| + |z_n|. \end{aligned}$$

**Exercise 1.13** If  $x, y$  are complex, prove that

$$||x| - |y|| \leq |x - y|.$$

*Solution.* Since  $x = x - y + y$ , the triangle inequality gives

$$|x| \leq |x - y| + |y|,$$

so that  $|x| - |y| \leq |x - y|$ . Similarly  $|y| - |x| \leq |x - y|$ . Since  $|x| - |y|$  is a real number we have either  $||x| - |y|| = |x| - |y|$  or  $||x| - |y|| = |y| - |x|$ . In either case, we have shown that  $||x| - |y|| \leq |x - y|$ .

**Exercise 1.14** If  $z$  is a complex number such that  $|z| = 1$ , that is, such that  $z\bar{z} = 1$ , compute

$$|1 + z|^2 + |1 - z|^2.$$

*Solution.*  $|1 + z|^2 = (1 + z)(1 + \bar{z}) = 1 + \bar{z} + z + z\bar{z} = 2 + z + \bar{z}$ . Similarly  $|1 - z|^2 = (1 - z)(1 - \bar{z}) = 1 - z - \bar{z} + z\bar{z} = 2 - z - \bar{z}$ . Hence

$$|1 + z|^2 + |1 - z|^2 = 4.$$

**Exercise 1.15** Under what conditions does equality hold in the Schwarz inequality?

*Solution.* The proof of Theorem 1.35 shows that equality can hold if  $B = 0$  or if  $Ba_j - Cb_j = 0$  for all  $j$ , i.e., the numbers  $a_j$  are proportional to the numbers  $b_j$ . (In terms of linear algebra this means the vectors  $\mathbf{a} = (a_1, a_2, \dots, a_n)$  and  $\mathbf{b} = (b_1, b_2, \dots, b_n)$  in complex  $n$ -dimensional space are linearly dependent. Conversely, if these vectors are linearly independent, then strict inequality holds.)

**Exercise 1.16** Suppose  $k \geq 3$ ,  $\mathbf{x}, \mathbf{y} \in R^k$ ,  $|\mathbf{x} - \mathbf{y}| = d > 0$ , and  $r > 0$ . Prove:

(a) If  $2r > d$ , there are infinitely many  $\mathbf{z} \in R^k$  such that

$$|\mathbf{z} - \mathbf{x}| = |\mathbf{z} - \mathbf{y}| = r.$$

(b) If  $2r = d$ , there is exactly one such  $\mathbf{z}$ .

(c) If  $2r < d$ , there is no such  $\mathbf{z}$ .

How must these statements be modified if  $k$  is 2 or 1?

*Solution.* (a) Let  $\mathbf{w}$  be any vector satisfying the following two equations:

$$\begin{aligned} \mathbf{w} \cdot (\mathbf{x} - \mathbf{y}) &= 0, \\ |\mathbf{w}|^2 &= r^2 - \frac{d^2}{4}. \end{aligned}$$

From linear algebra it is known that all but one of the components of a solution  $\mathbf{w}$  of the first equation can be arbitrary. The remaining component is then uniquely determined. Also, if  $\mathbf{w}$  is any non-zero solution of the first equation, there is a unique positive number  $t$  such that  $t\mathbf{w}$  satisfies both equations. (For example, if  $x_1 \neq y_1$ , the first equation is satisfied whenever

$$z_1 = \frac{z_2(x_2 - y_2) + \dots + z_k(x_k - y_k)}{y_1 - x_1}.$$

If  $(z_1, z_2, \dots, z_k)$  satisfies this equation, so does  $(tz_1, tz_2, \dots, tz_k)$  for any real number  $t$ .) Since at least two of these components can vary independently, we can find a solution with these components having any prescribed ratio. This



ratio does not change when we multiply by the positive number  $t$  to obtain a solution of both equations. Since there are infinitely many ratios, there are infinitely many distinct solutions. For each such solution  $\mathbf{w}$  the vector  $\mathbf{z} = \frac{1}{2}\mathbf{x} + \frac{1}{2}\mathbf{y} + \mathbf{w}$  is a solution of the required equation. For

$$\begin{aligned} |\mathbf{z} - \mathbf{x}|^2 &= \left| \frac{\mathbf{y} - \mathbf{x}}{2} + \mathbf{w} \right|^2 \\ &= \left| \frac{\mathbf{y} - \mathbf{x}}{2} \right|^2 + 2\mathbf{w} \cdot \frac{\mathbf{x} - \mathbf{y}}{2} + |\mathbf{w}|^2 \\ &= \frac{d^2}{4} + 0 + r^2 - \frac{d^2}{4} \\ &= r^2, \end{aligned}$$

and a similar relation holds for  $|\mathbf{z} - \mathbf{y}|^2$ .

(b) The proof of the triangle inequality shows that equality can hold in this inequality only if it holds in the Schwarz inequality, i.e., one of the two vectors is a scalar multiple of the other. Further examination of the proof shows that the scalar must be nonnegative. Now the conditions of this part of the problem show that

$$|\mathbf{x} - \mathbf{y}| = d = |\mathbf{x} - \mathbf{z}| + |\mathbf{z} - \mathbf{y}|.$$

Hence it follows that there is a nonnegative scalar  $t$  such that

$$\mathbf{x} - \mathbf{z} = t(\mathbf{z} - \mathbf{y}).$$

However, the hypothesis also shows immediately that  $t = 1$ , and so  $\mathbf{z}$  is uniquely determined as

$$\mathbf{z} = \frac{\mathbf{x} + \mathbf{y}}{2}.$$

(c) If  $\mathbf{z}$  were to satisfy this condition, the triangle inequality would be violated, i.e., we would have

$$|\mathbf{x} - \mathbf{y}| = d > 2r = |\mathbf{x} - \mathbf{z}| + |\mathbf{z} - \mathbf{y}|.$$

When  $k = 2$ , there are precisely 2 solutions in case (a). When  $k = 1$ , there are no solutions in case (a). The conclusions in cases (b) and (c) do not require modification.

**Exercise 1.17** Prove that

$$|\mathbf{x} + \mathbf{y}|^2 + |\mathbf{x} - \mathbf{y}|^2 = 2|\mathbf{x}|^2 + 2|\mathbf{y}|^2$$

if  $\mathbf{x} \in R^k$  and  $\mathbf{y} \in R^k$ . Interpret this geometrically as a statement about parallelograms.

*Solution.* The proof is a routine computation, using the relation

$$|\mathbf{x} \pm \mathbf{y}|^2 = (\mathbf{x} \pm \mathbf{y}) \cdot (\mathbf{x} \pm \mathbf{y}) = |\mathbf{x}|^2 \pm 2\mathbf{x} \cdot \mathbf{y} + |\mathbf{y}|^2.$$

If  $\mathbf{x}$  and  $\mathbf{y}$  are the sides of a parallelogram, then  $\mathbf{x} + \mathbf{y}$  and  $\mathbf{x} - \mathbf{y}$  are its diagonals. Hence this result says that the sum of the squares on the diagonals of a parallelogram equals the sum of the squares on the sides.

**Exercise 1.18** If  $k \geq 2$  and  $\mathbf{x} \in R^k$ , prove that there exists  $\mathbf{y} \in R^k$  such that  $\mathbf{y} \neq \mathbf{0}$  but  $\mathbf{x} \cdot \mathbf{y} = 0$ . Is this also true if  $k = 1$ ?

*Solution.* If  $\mathbf{x}$  has any components equal to 0, then  $\mathbf{y}$  can be taken to have the corresponding components equal to 1 and all others equal to 0. If all the components of  $\mathbf{x}$  are nonzero,  $\mathbf{y}$  can be taken as  $(-x_2, x_1, 0, \dots, 0)$ . This is, of course, not true when  $k = 1$ , since the product of two nonzero real numbers is nonzero.

**Exercise 1.19** Suppose  $\mathbf{a} \in R^k$ ,  $\mathbf{b} \in R^k$ . Find  $\mathbf{c} \in R^k$  and  $r > 0$  such that

$$|\mathbf{x} - \mathbf{a}| = 2|\mathbf{x} - \mathbf{b}|$$

if and only if  $|\mathbf{x} - \mathbf{c}| = r$ . (*Solution:*  $3\mathbf{c} = 4\mathbf{b} - \mathbf{a}$ ,  $3r = 2|\mathbf{b} - \mathbf{a}|$ .)

*Solution.* Since the solution is given to us, all we have to do is verify it, i.e., we need to show that the equation

$$|\mathbf{x} - \mathbf{a}| = 2|\mathbf{x} - \mathbf{b}|$$

is equivalent to  $|\mathbf{x} - \mathbf{c}| = r$ , which says

$$\left| \mathbf{x} - \frac{4}{3}\mathbf{b} + \frac{1}{3}\mathbf{a} \right| = \frac{2}{3}|\mathbf{b} - \mathbf{a}|.$$

If we square both sides of both equations, we an equivalent pair of equations, the first of which reduces to

$$3|\mathbf{x}|^2 + 2\mathbf{a} \cdot \mathbf{x} - 8\mathbf{b} \cdot \mathbf{x} - |\mathbf{a}|^2 + 4|\mathbf{b}|^2 = 0,$$

and the second of which reduces to this equation divided by 3. Hence these equations are indeed equivalent.

**Exercise 1.20** With reference to the Appendix, suppose that property (III) were omitted from the definition of a cut. Keep the same definitions of order and addition. Show that the resulting ordered set has the least-upper-bound property, that addition satisfies axioms (A1) to (A4) (with a slightly different zero element!) but that (A5) fails.

*Solution.* We are now defining a cut to be a proper subset of the rational numbers that contains, along with each of its elements, all smaller rational

numbers. Order is defined by containment. Now given a set  $A$  of cuts having an upper bound  $\beta$ , let  $\alpha$  be the union of all the cuts in  $A$ . Obviously  $\alpha$  is properly contained in  $\beta$ , and so is a proper subset of the rationals. It also obviously satisfies the property that if  $p \in \alpha$  and  $q < p$ , then  $q \in \alpha$ ; hence  $\alpha$  is a cut. It is further obvious that  $\alpha$  contains each elements of  $A$ , and so is an upper bound for  $A$ . It remains to prove that there is no smaller upper bound.

To that end, suppose,  $\gamma < \alpha$ , then  $\alpha$  contains an element  $x$  not in  $\gamma$ . By definition of  $\alpha$ ,  $x$  must belong to some cut  $\delta$  in  $A$ . But then  $\gamma < \delta$ , and so  $\gamma$  is not an upper bound for  $A$ . Thus  $\alpha$  is the least upper bound.

The proof given in the text goes over without any change to show that (A1), (A2), and (A3) hold. As for (A4) let  $O = \{r : r \leq 0\}$ . We claim  $O + \alpha = \alpha$ . The proof is easy. First, we obviously have  $O + \alpha \subseteq \alpha$ . For  $r + s \leq s$  if  $r \leq 0$ . Hence  $r + s \in \alpha$  if  $s \in \alpha$ . Conversely  $\alpha \subseteq O + \alpha$ , since each  $s$  in  $\alpha$  can be written as  $0 + s$ .

Unfortunately, if  $O' = \{r : r < 0\}$ , there is no element  $\alpha$  such that  $\alpha + O' = O$ . For  $\alpha + O'$  has no largest element. If  $x = r + s \in \alpha + O'$ , where  $r \in \alpha$  and  $s \in O'$ , there is an element  $t \in O'$  with  $t > s$ , and so  $r + t \in \alpha + O'$  and  $r + t > s$ . Since  $O$  has a largest element (namely 0), these two sets cannot be equal.