

Updated Errata for “CS402 - Theory of Automata” Lecture Notes

Dear students,

If you have an older copy of cs402 handouts then there is possibility of below mentioned errors in them please change them accordingly in your printed copy. (Changes have been highlighted in bold font)

| | |
|--------------------------------------|--|
| Lecture No: 2 | |
| Page No: 8 | |
| Old Text | New / Updated Text |
| N / A (Add it at the end of page 8.) | <p>Defining the language L, of strings containing exactly two a’s, defined over $\Sigma=\{a, b\}$</p> <p>Step 1: a is in L</p> <p>Step 2: s(aa)s is also in L, where s belongs to b^*</p> <p>Step 3: No strings except those constructed in above, are allowed to be in L</p> |

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|---|--|
| Lecture No: 8 | |
| Page No: 20 | |
| Old Text | New / Updated Text |
| <p>Example</p> <p>Consider the language L of strings, defined over $\Sigma=\{a, b\}$, not ending in b. The language L may be expressed by RE $(a + b)^* a$, may be accepted by the following TG</p> <p>(Given TG is not correct..)</p> | <p>Correct TG is as below;</p> <pre> graph TD Start((-)) -- "a" --> Plus((+)) Start -- "Λ, a" --> Plus Plus -- "a" --> Plus Plus -- "a, b" --> Plus style Plus stroke-width:4px </pre> |

Errata for “Theory of Automata” Lecture Notes

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Lecture No. 1

Page No. 4 Topic: Introduction to Computer Theory

| Old Text | Change |
|---|---|
| Example: $\Sigma = \{B, aB, bab, d\}$ $s = BaBbabBd$ Tokenizing = (B), (aB), (bab), (d) $ s = 4$ | Example: $\Sigma = \{B, aB, bab, d\}$ $s = BaBbabBd$ Tokenizing = (B), (aB), (bab), (B), (d) $ s = 5$ |

Lecture No. 2

Page No 7 Topic: Introduction to Computer Theory

| Old Text | Change |
|--|--|
| Example: The language EQUAL , of strings with number of a's equal to number of b's, defined over $\Sigma = \{a, b\}$, can be written as $\{\Lambda, ab, aabb, abab, baba, abba, \dots\}$ | Example: The language EQUAL , of strings with number of a's equal to number of b's, defined over $\Sigma = \{a, b\}$, can be written as $\{\Lambda, ab, ba, aabb, abab, baba, abba, \dots\}$ |

Page No 8 Topic: Introduction to Computer Theory

| Old Text | Change |
|--|---|
| <u>Defining the language L, of strings containing exactly one a, defined over $\Sigma = \{a, b\}$</u> Step 1: a is in L Step 2: s(aa)s is also in L, where s belongs to b^* Step 3: No strings except those constructed in above, are allowed to be in L | <u>Defining the language L, of strings containing exactly one a, defined over $\Sigma = \{a, b\}$</u> Step 1: a is in L Step 2: s(a)s is also in L, where s belongs to b^* Step 3: No strings except those constructed in above, are allowed to be in L |

Lecture No. 25

Page No. 68

Topic: Intersection of two regular languages

| Old Text | Change |
|---|---|
| Now FA accepting $L_1^c \cap L_2^c$, using the method described earlier, may be as follows | Now FA accepting $L_1^c \cup L_2^c$, using the method described earlier, may be as follows |

Lecture No. 35

Page No. 100 Topic: Nullable Production

| Old Text | Change |
|---|---|
| Consider the following CFG $S \rightarrow XaY YY aX ZYX$ $X \rightarrow Za bZ ZZ Yb$ $Y \rightarrow Ya XY \wedge$ $Z \rightarrow aX YYY$ It is to be noted that in the given CFG, the productions $S \rightarrow YY$, $X \rightarrow ZZ$, $Z \rightarrow YYY$ are Nullable productions, while $Y \rightarrow \wedge$ is Null production. | Consider the following CFG $S \rightarrow XaY YY aX ZYX$ $X \rightarrow Za bZ ZZ Yb$ $Y \rightarrow Ya XY \wedge$ $Z \rightarrow aX YYY$ It is to be noted that in the given CFG, the productions $S \rightarrow YY$, $S \rightarrow ZYX$, $X \rightarrow ZZ$, $Z \rightarrow YYY$ are Nullable productions, while $Y \rightarrow \wedge$ is Null production. |
| Here the method of removing null productions, as discussed earlier, will be used along with replacing | Here the method of removing null productions, as discussed earlier, will be used along with replacing nonterminals corresponding to nullable productions |

| | |
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| <p>nonterminals corresponding to nullable productions like nonterminals for null productions are replaced. Thus the required CFG will be $S \rightarrow XaY Xa aY a YY Y aX ZYX YX ZX ZY Z Y X$ $X \rightarrow Za a bZ b ZZ Z Yb$ $Y \rightarrow Ya a XY X Y$ $Z \rightarrow aX a YYY YY Y,$</p> | <p>like nonterminals for null productions are replaced. Thus the required CFG will be $S \rightarrow XaY Xa aY a YY Y aX ZYX YX ZX ZY Z Y X$ $X \rightarrow Za a bZ b ZZ Z Yb$ $Y \rightarrow Ya a XY X Y$ $Z \rightarrow aX a YYY YY Y,$</p> |
|--|---|

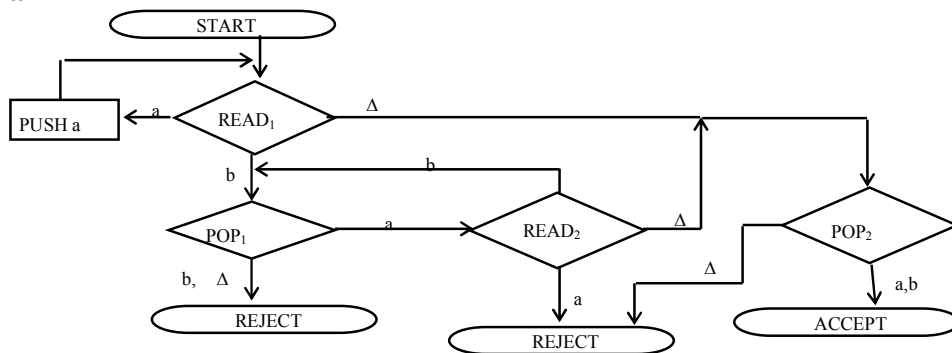
Lecture No. 37

Page No. 107

Topic: **POP and STACK**

Change (20-12-2011)

Old Diagram



Corrected

